

**Orientable
quadrilateral embeddings
of cartesian products**

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Quadrilateral embeddings and cartesian products

Orientable surfaces:

S_h



Embedding $\Phi : G \rightarrow \Sigma$: draw G in Σ without edge crossings.

Quadrilateral: open disk face bounded by 4-cycle.

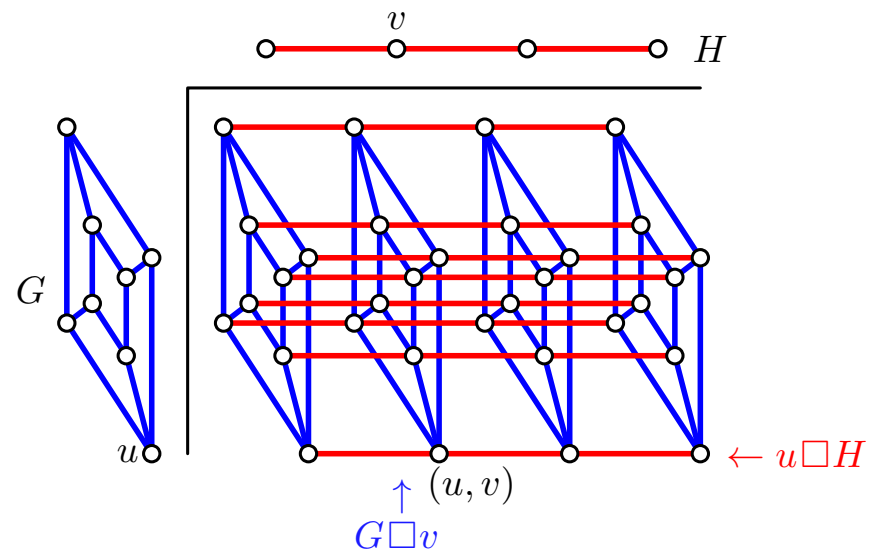
Quadrilateral embedding (QE): every face quadrilateral. So cellular.

Why quadrilateral embeddings? Minimum genus if graph has girth 4 or more.

Cartesian product (CP) $G \square H$:

G -edges inside $G \square v$, H -edges
inside $u \square H$.

Why cartesian products? Many
4-cycles, improves chances of
finding quadrilateral embedding.



Pisanski's three questions, 1992

Question 1: If G, H are arbitrary 1-factorable t -regular graphs, does $G \square H$ always have an orientable quadrilateral embedding?

True if G, H bipartite (Pisanski, 1980).

Question 2: For t -regular G , $t \geq 2$, does

$$G \square C_{2n_1} \square C_{2n_2} \square \dots \square C_{2n_t}$$

have an orientable quadrilateral embedding?

More general than $G \square Q_{2t} = G \square (\square^t C_4)$.

True if G bipartite (Pisanski, 1980).

Question 3: For an arbitrary graph G , does $G \square Q_n$ have an orientable quadrilateral embedding for all sufficiently large n ? ($Q_n = \square^n K_2$, hypercube.)

True if G bipartite, for $n \geq \Delta(G)$ (Pisanski, 1980).

True for regular G if $n \geq 2\Delta(G) + 3$ (Pisanski, 1992).

Also true for all G if we drop 'orientable' (Pisanski, 1982 and also Hunter and Kainen, 2007).

Here we discuss Questions 2 and 3 ...

Our construction

Generalizes Pisanski's +/- construction, 1992.

Pisanski showed that for every t -regular G , there is an orientable QE of $G \square Q_n$ for all $n \geq 2t + 3$.

- Begin with orientable emb. Φ of any graph G .

Add **semiedges** coloured by D , $|D| = r$: Φ^+ where

(0) each colour appears once at each vertex,

(1) edge/semiedge adjacency condition

(\rightarrow GH -faces),

(2) faces without semiedges are quadrilaterals

(\rightarrow G -faces).

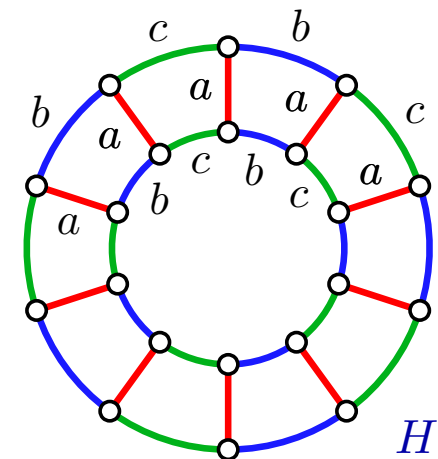
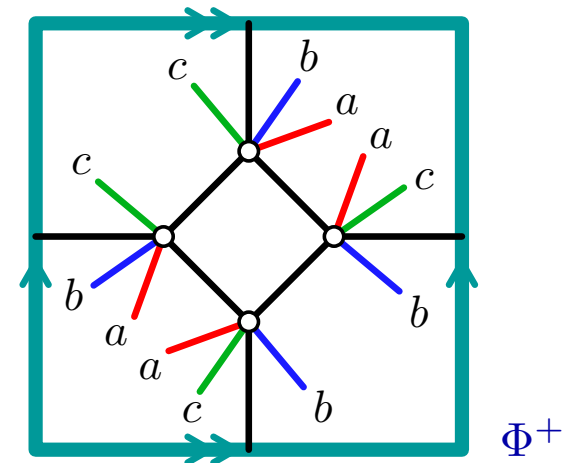
- Colour edges of r -regular bipartite H with D so

(3) consecutive colours d_1, d_2 in Φ^+ mean

$H(d_1, d_2)$ is a 4-cycle 2-factor (\rightarrow H -faces).

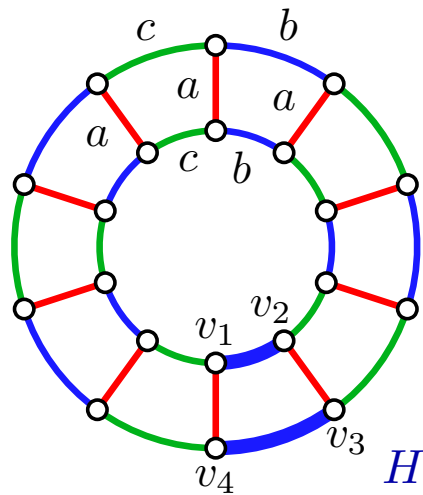
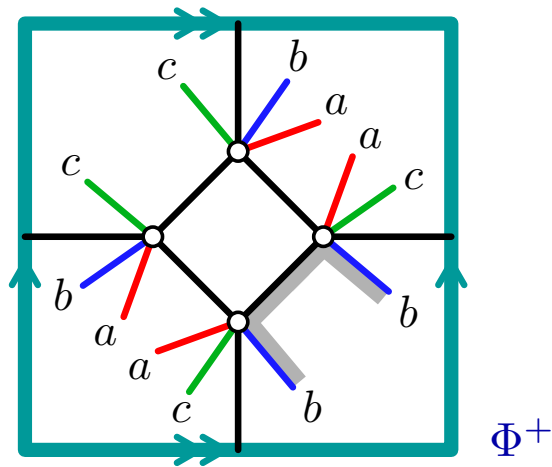
- Use to derive orientable QE of $G \square H$.

Example: $K_4 \square (C_{10} \square K_2)$

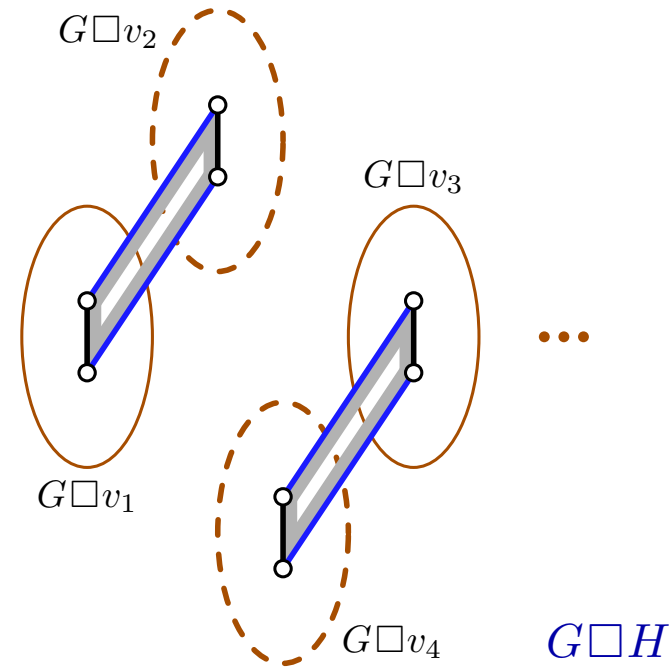


Construction details I

Example: $K_4 \square (C_{10} \square K_2)$

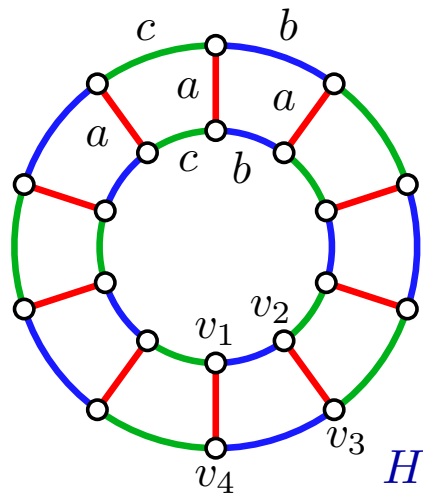
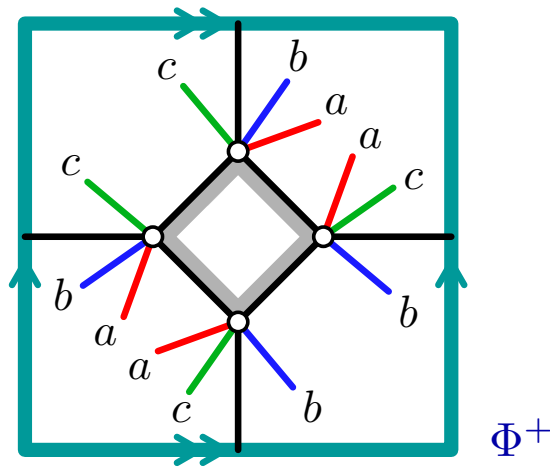


(1) Get GH -faces from corners between edges and semiedges, using edge/semiedge adjacency condition.

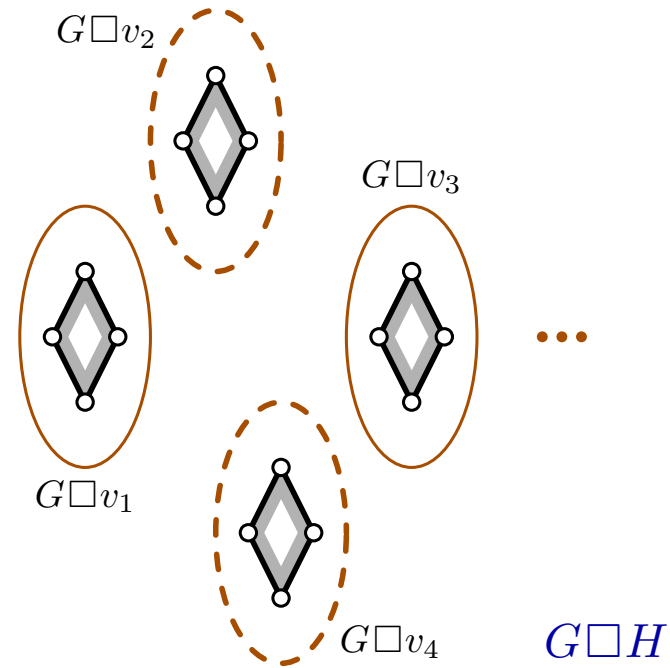


Construction details II

Example: $K_4 \square (C_{10} \square K_2)$

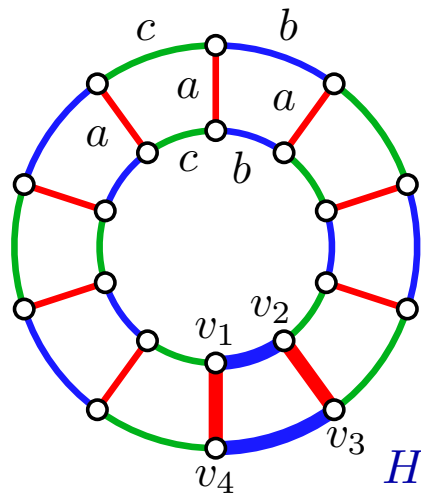
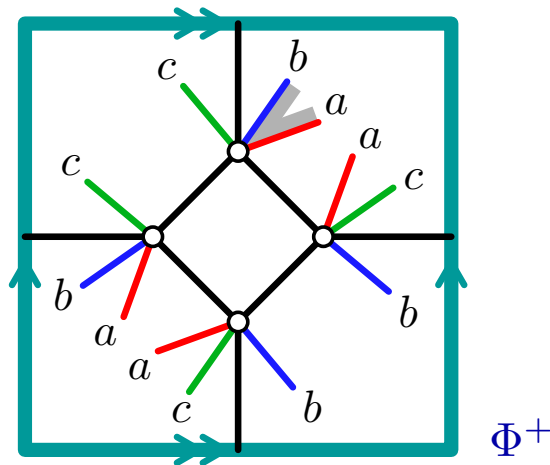


(2) Get G -faces from corners between pairs of edges, using fact that faces without semiedges are quadrilaterals.

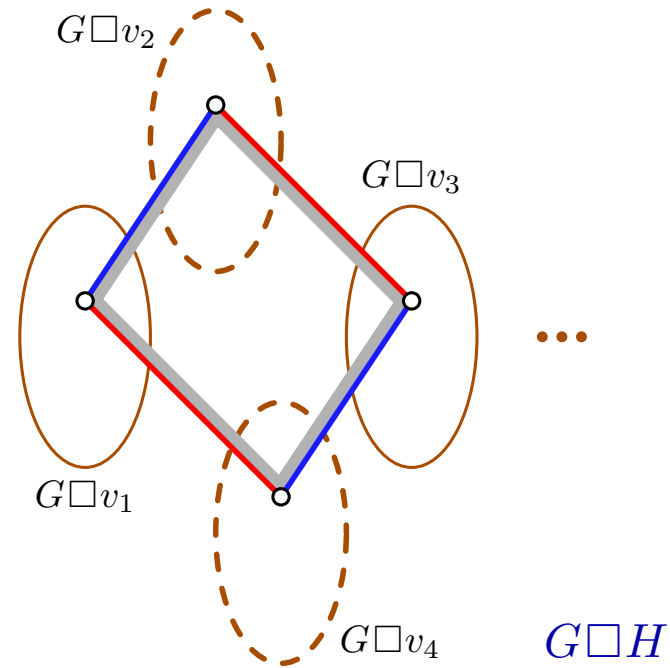


Construction details III

Example: $K_4 \square (C_{10} \square K_2)$



(3) Get H -faces from corners between pairs of semiedges, using fact that consecutive colours d_1, d_2 in Φ^+ mean $H(d_1, d_2)$ is a 4-cycle 2-factor.

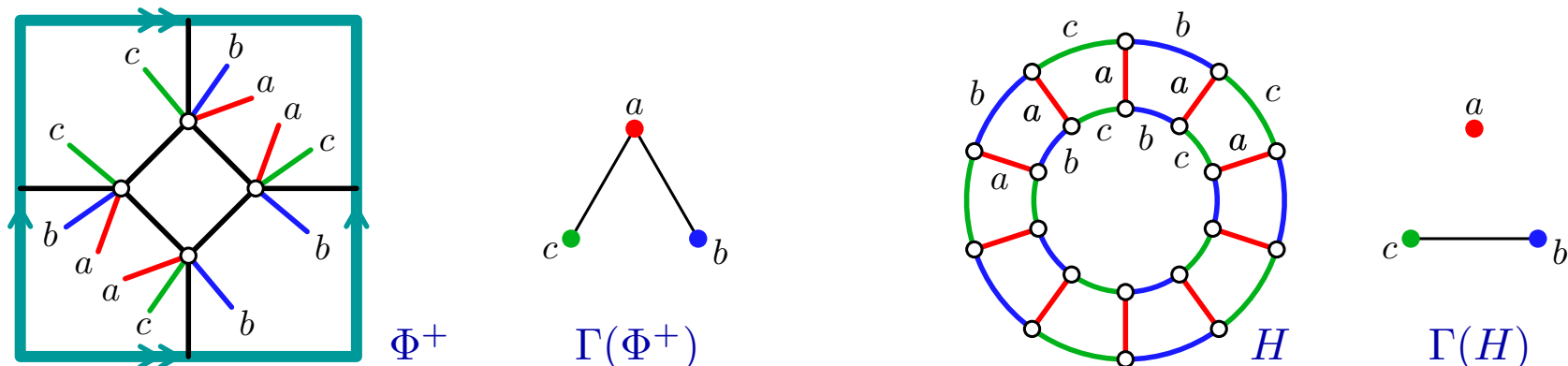


Conflict graphs

Hardest part is satisfying (3). Think of Φ^+ and H as generating **conflicts** between pairs of colours:

- conflict in Φ^+ if d_1, d_2 consecutive somewhere,
- conflict in H if $H(d_1, d_2)$ not a 4-cycle 2-factor.

Want **conflict graphs** $\Gamma(\Phi^+)$, $\Gamma(H)$ to be edge-disjoint. For example:



- Can weaken this. Enough if $\Gamma(\Phi^+)$ and $\Gamma(H)$ **pack**: one isomorphic to subgraph of complement of other. Can also use different colours for Φ^+ , H .
- If H is itself a cartesian product of regular graphs all of the same degree (e.g., H is a cube) then we can use **equitable colourings** of $\Gamma(\Phi^+)$ to show that $\Gamma(\Phi^+)$ and $\Gamma(H)$ **pack**: **Hajnal-Szemerédi Theorem** or special construction.

Solving Questions 2 and 3

From equitable colourings we get:

Theorem: Suppose that G is k -edge-colourable, $k \geq 3$, and H_1, H_2, \dots, H_m are all s -regular bipartite graphs, where $m \geq 3$ and $sm \geq \lceil 3k/2 \rceil$. Then $G \square (H_1 \square H_2 \square \dots \square H_m)$ has an orientable quadrilateral embedding.

Question 2: For t -regular G , $t \geq 2$, does

$$G \square C_{2n_1} \square C_{2n_2} \square \dots \square C_{2n_t}$$

have an orientable quadrilateral embedding?

Answer: Yes, for $t \geq 3$. In fact, works for

$$G \square C_{2n_1} \square C_{2n_2} \square \dots \square C_{2n_m}$$

provided $t \geq 2$ and $m \geq \max(3, \lceil 3(t+1)/4 \rceil)$.

Question 3: For an arbitrary graph G , does $G \square Q_n$ have an orientable quadrilateral embedding for all sufficiently large n ?

Answer: Yes. Just take all $H_i = K_2$, then $n \geq \max(3, \lceil 3\chi'(G)/2 \rceil)$ works. ($\chi'(G)$, chromatic index, is $\Delta(G)$ or $\Delta(G) + 1$.)

Future directions

- Extend our construction for $G \square H$:
 - Nonorientable embeddings: start with nonorientable embedding of G , or use nonbipartite H .
 - Nonregular graphs H , using partial 4-cycle patchworks, or directly.
- What about orientable quadrilateral embeddings of $G \square H$ when neither G nor H is bipartite? **Nothing much known.**
- We have 3-regular counterexamples to Question 1: no orientable QE of $G \square H$ for G, H both t -regular, 1-factorable. Find counterexamples for Question 1 that are t -regular for $t \geq 4$. **Should be doable.**
- What about Question 2 for 2-regular G ? Does $C_{\text{odd}} \square C_{\text{even}} \square C_{\text{even}}$ have an orientable quadrilateral embedding? **Our technique does not work.**

Thank you!

And congratulations to Brian and Dragan!