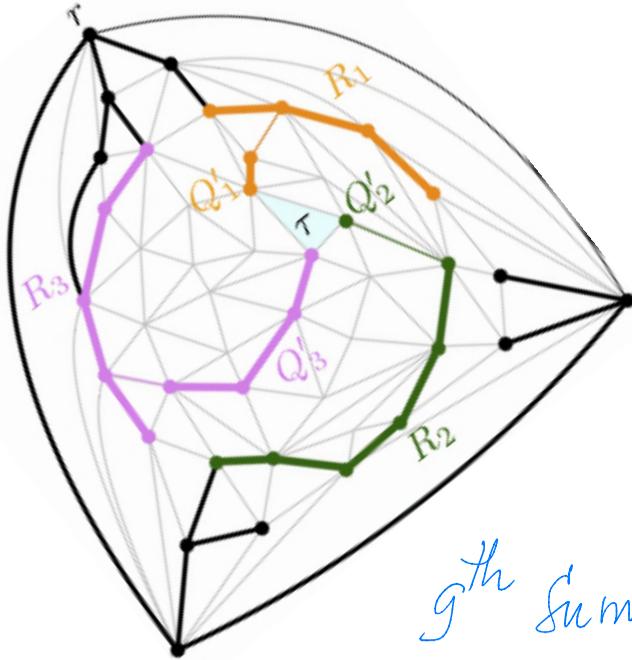


# layered partitions with applications

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# APPLICATIONS

- queue layouts [1992]  
& 3D grid drawings [2002]
- non-repetitive graph colourings [2002]
- $p$ -centered colourings
- clustered colourings
- $\vdots$
- simple proofs of known results, graph products, ...

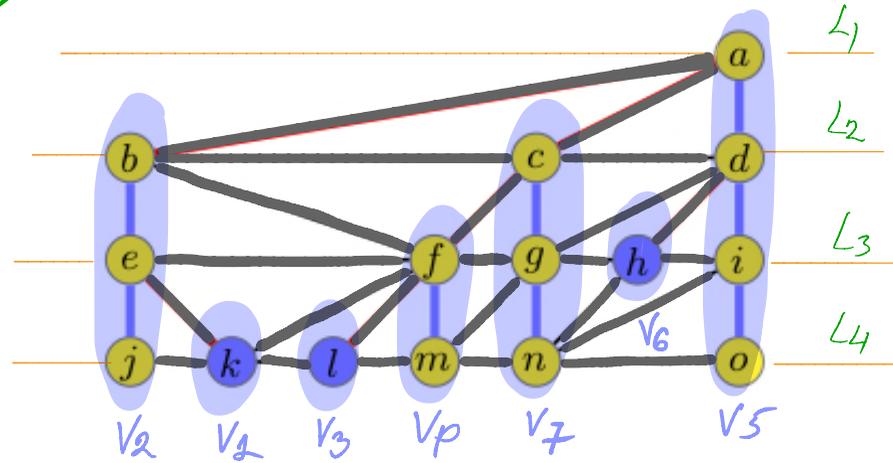
simple proofs  
+ better bounds

# tool: LAYERED PARTITIONS

layering  $\mathcal{L} = \langle L_1, L_2, \dots \rangle$  of  $G$

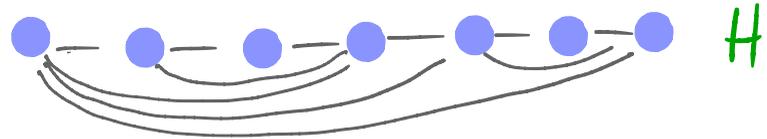
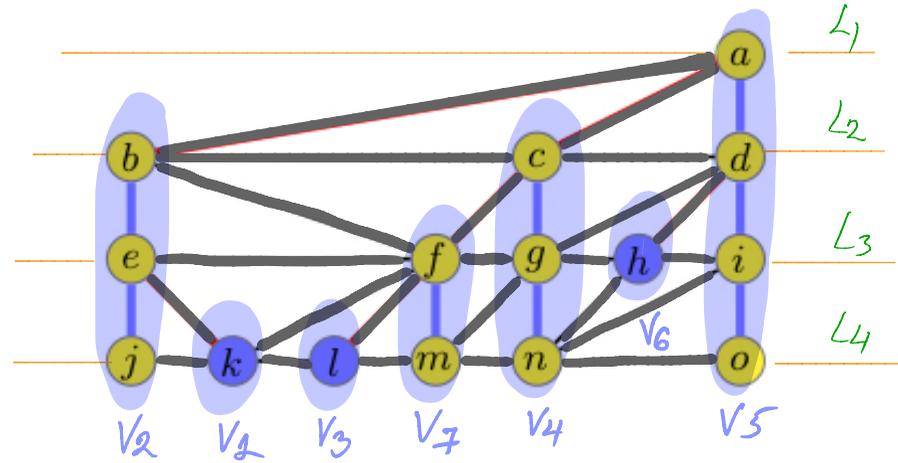
partition  $\mathcal{P} = \{V_1, \dots, V_p\}$  of  $V(G)$

layered partition:  $\Rightarrow$   
a pair  $(\mathcal{L}, \mathcal{P})$



# tool: LAYERED PARTITIONS

layered partition:  
a pair  $(\mathcal{L}, \mathcal{P})$



layered width:

$$\max \{ |L \cap P| : L \in \mathcal{L}, P \in \mathcal{P} \}$$

quotient graph  $H = G/P$

WANT

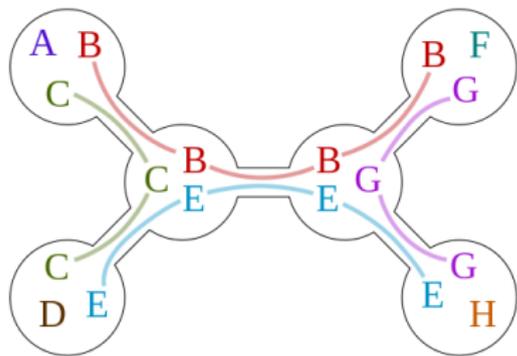
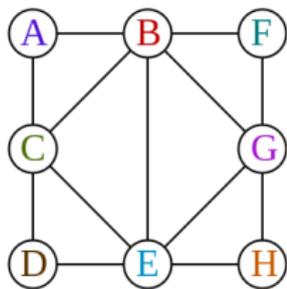
layered  $H$ -partitions of  $G$  with :

small (constant) layered width

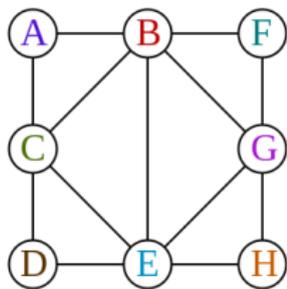
+  
nice quotient graph  $H$

↳ bounded tree width

a **tree decomposition** represents each vertex as a subtree of a tree  $T$  so that the subtrees of adjacent vertices intersect in  $T$

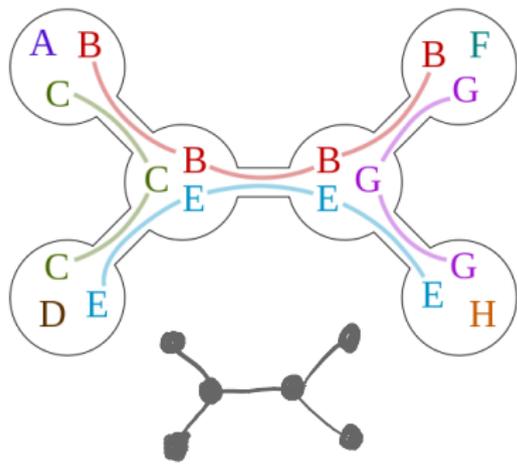


a **tree decomposition** represents each vertex as a subtree of a tree  $T$  so that the subtrees of adjacent vertices intersect in  $T$



**tree-width** := maximum bag size - 1

↳ measure of how 'tree-like' a graph is



•  $T$

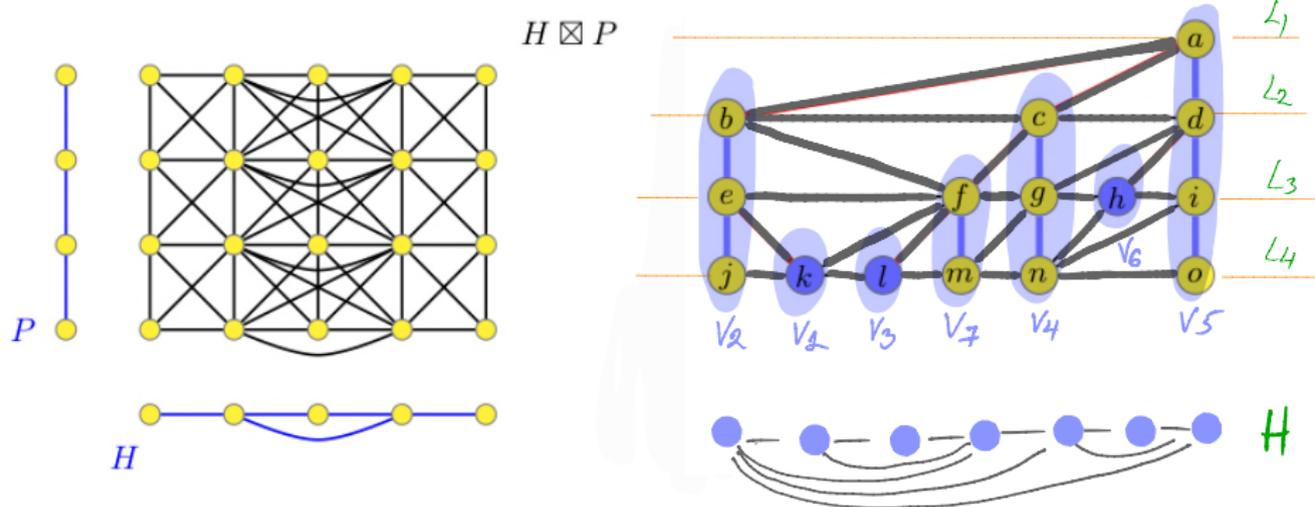
## which graph classes have it?

- apex minor-free graphs  
(planar, bounded genus, ...) } minor closed  
[D., Joret, Micc, Morin, Ueckerdt, Wood] 2019
- $k$ -planar,  $d$ -map graphs ... } non-minor closed  
[D., Morin, Wood] 2019

# structure of planar graphs

Corollary

every planar graph  $G$  is the subgraph of  $H \boxtimes P$   
for some graph  $H$  with treewidth  $\leq 8$  and some path  $P$



First there was...

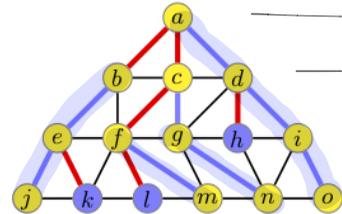
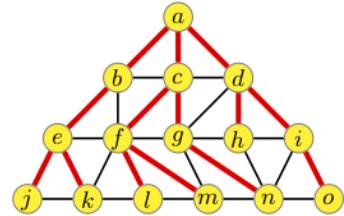
Th: [Mi. Pilipczuk, Siebertz] 2018  
Every planar graph  $G$  has a partition  $\mathcal{P}$   
into geodesics s.t.  $G/\mathcal{P}$  has treewidth  $\leq 8$ .

# Main Theorem

th: [D., Joret, Micoc, Morin, Ueckerdt, Wood] 2019

For every BFS spanning tree  $T$  of a planar graph  $G$ ,  
 $\exists$  a partition  $\mathcal{P}$  into vertical paths s.t.  $\text{tree width}(G/\mathcal{P}) \leq 8$ .

- gives layered  $H$ -partitions  
of layered width 1 &  $\text{tw}(H) \leq 8$



## Main Theorem: Proof

th: [D., Joret, Micoc, Morin, Ueckerdt, Wood] 2019

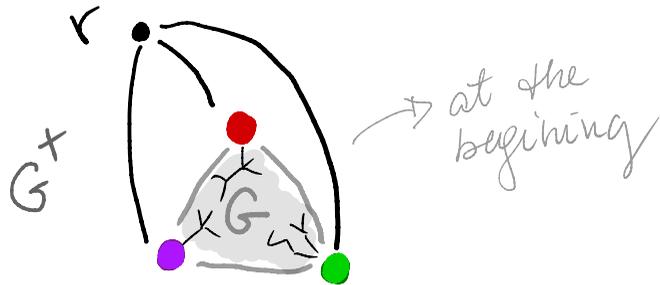
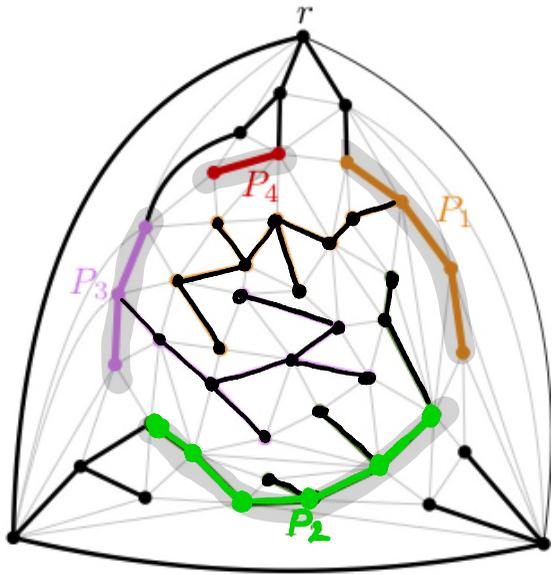
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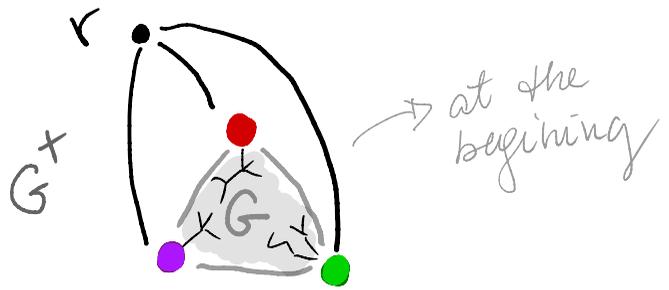
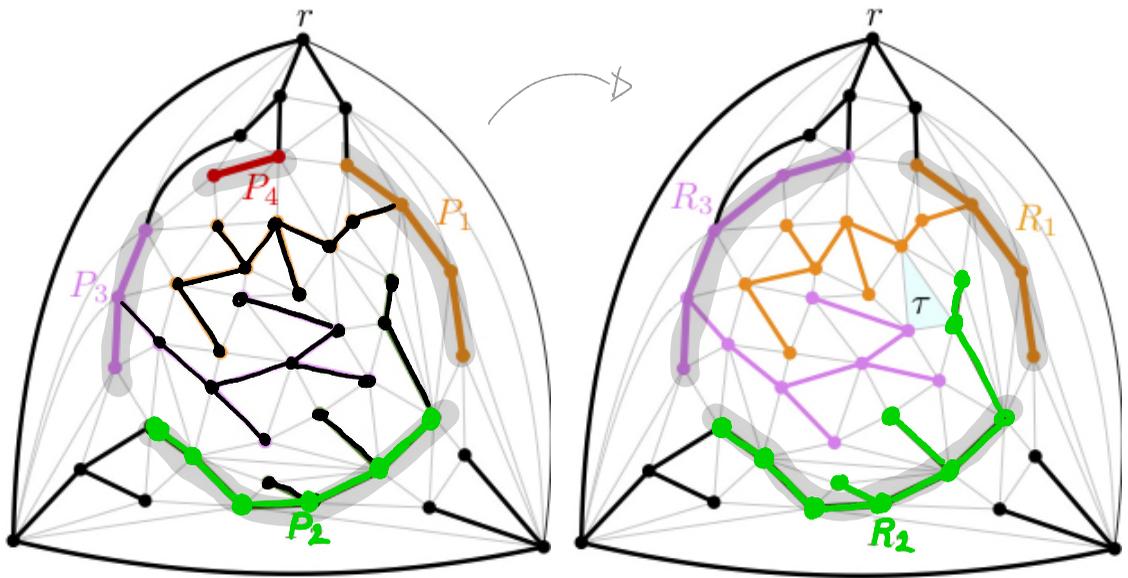
# Proof: Partitioning Planar Graphs

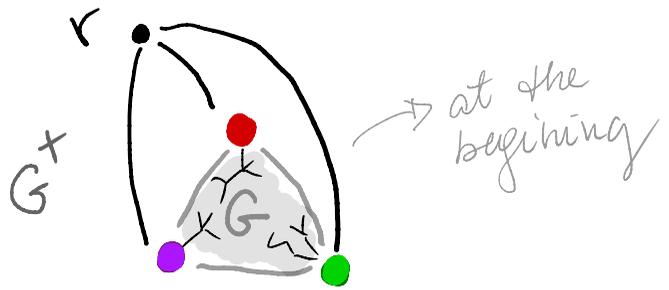
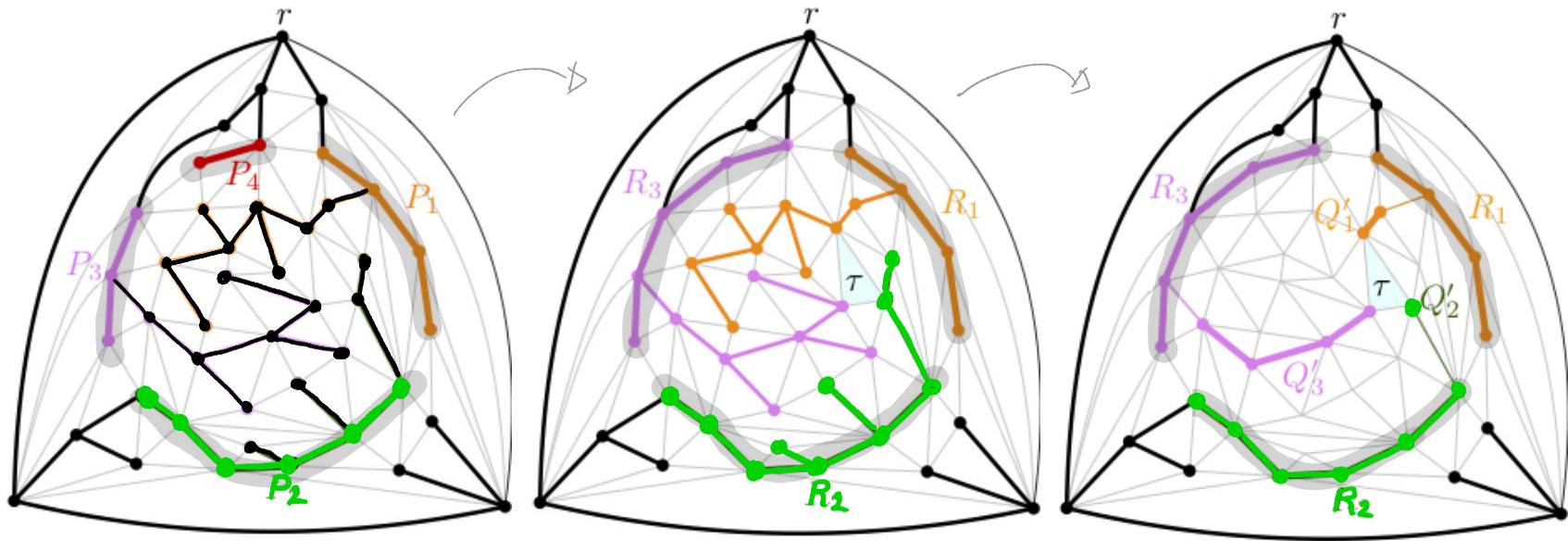
**Key lemma.** Suppose

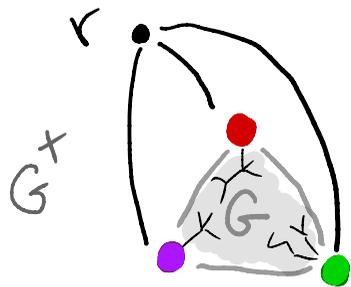
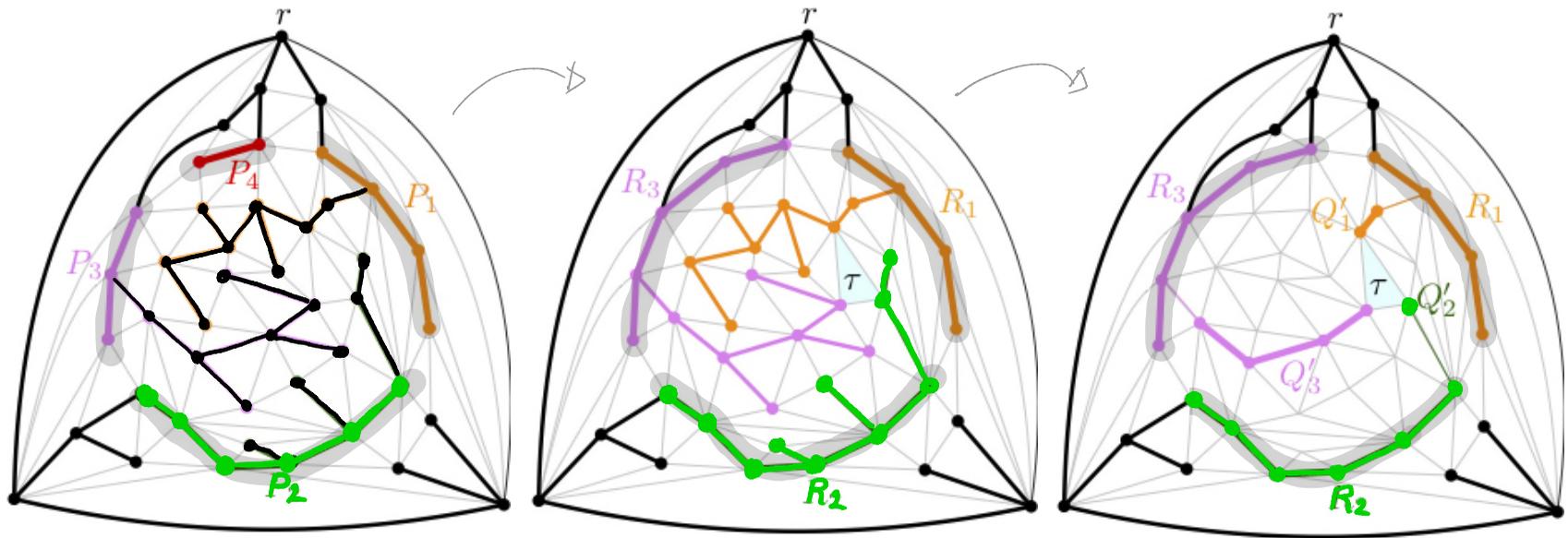
- $G^+$  plane triangulation
- $T$  rooted spanning tree of  $G^+$  with root on outer-face
- cycle  $C$  partitioned into vertical paths  $P_1, \dots, P_k$ , with  $k \leq 6$
- $G$  near-triangulation consisting of  $C$  and everything inside.

Then  $G$  has a partition  $\mathcal{P}$  into vertical paths where  $P_1, \dots, P_k \in \mathcal{P}$   
s.t.  $G/\mathcal{P}$  has a tree-decomposition in which every bag has size at most 9 and some bag contains all vertices corresponding to  $P_1, \dots, P_k$ .

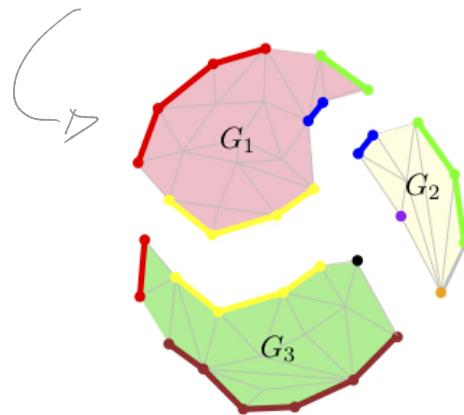








→ at the beginning



# Applications

**Dujmović, J., Micek, Morin, Ueckerdt, Wood '19** Planar graphs have bounded queue-number

**Dujmović, Esperet, J., Walczak, Wood '19** Planar graphs have bounded nonrepetitive chromatic number

*k*-queue layout of graph  $G$ :

- vertex ordering  $v_1, \dots, v_n$  of  $G$
- partition  $E_1, \dots, E_k$  of  $E(G)$  such that no two edges in  $E_i$  are nested



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**queue-number**  $qn(G) :=$  minimum  $k$  such that  $G$  has a  $k$ -queue layout

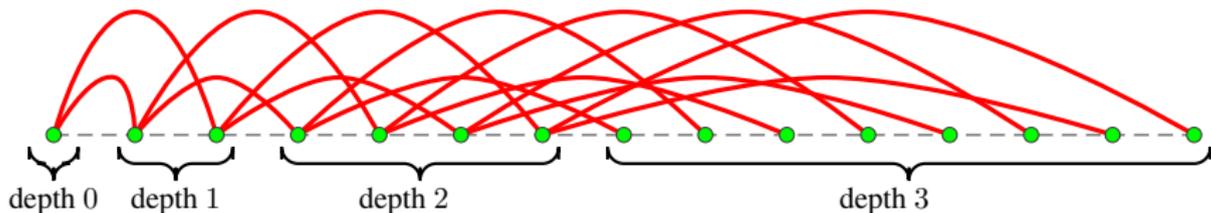
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**example**  $qn(\text{tree})=1$



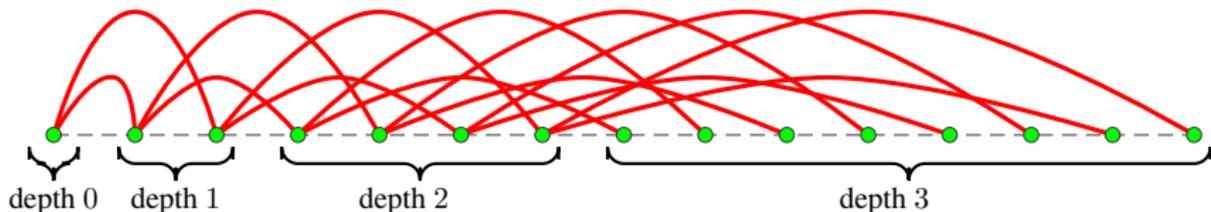
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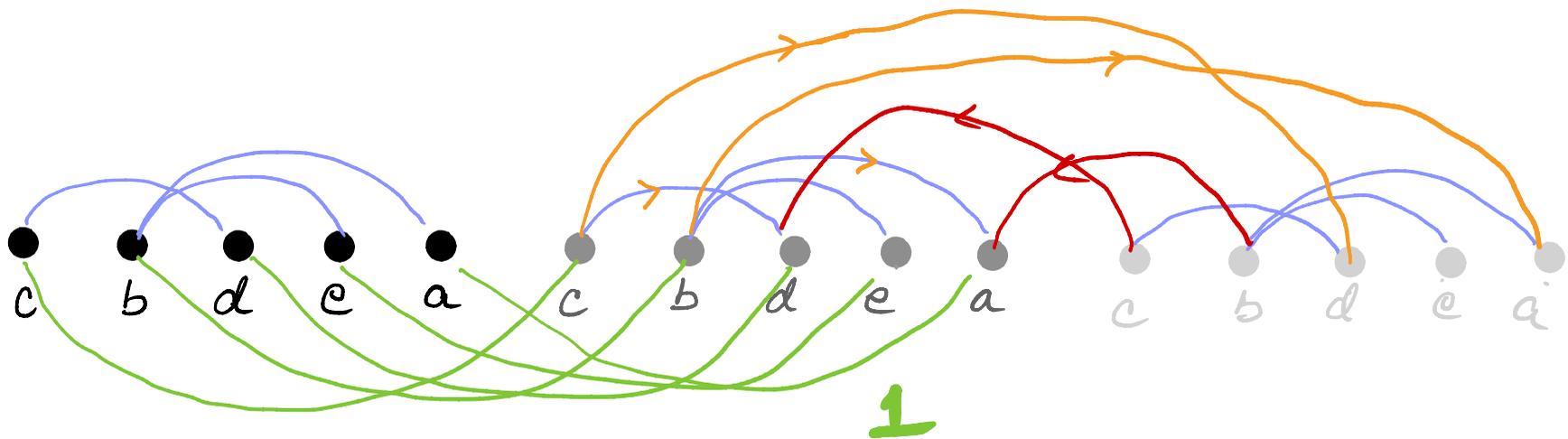
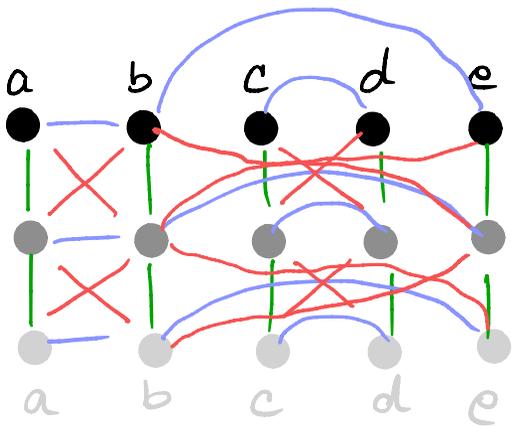
**open problem** [Heath, Leighton, Rosenberg '92]  
do planar graphs have bounded queue-number?

Using layered  $H$ -partitions of layered width  $l$

lemma: [DJMMUW '19 & Wood '08]

$$g_n(G) \leq 3 \cdot l \cdot g_n(H) + \lfloor \frac{3}{2} l \rfloor$$

Proof by picture



## Using layered $H$ -partitions of layered width $l$

Lemma: [DJMMUW '19 & Wood '08]

$$gn(G) \leq 3 \cdot l \cdot gn(H) + \lfloor \frac{3}{2} l \rfloor$$

KNOWN:

- bounded treewidth graphs have bounded queue-number [D., Morin, Wood '08]
- $gn(H) \leq 2^{tw(H)} - 1$  [Wiechert '18]

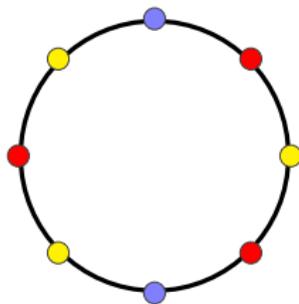
Theorem:

$$gn(\text{planar } G) \leq 3(2^8 - 1) + 1 = 766$$

[D., Joret, Miccic, Morin, Ueckerdt, Wood] 2019

## nonrepetitive colourings

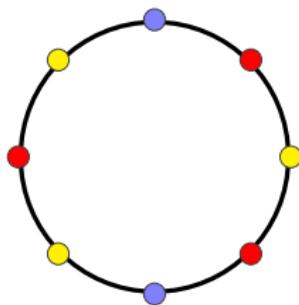
a colouring of a path  $v_1, \dots, v_{2t}$  is **repetitive**  
if  $\text{col}(v_i) = \text{col}(v_{t+i})$  for  $i \in \{1, \dots, t\}$



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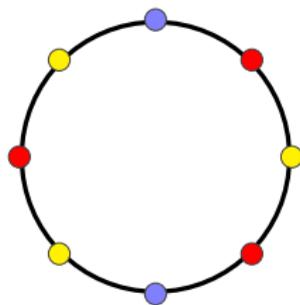
a colouring of a graph is **nonrepetitive**  
if no subpath is repetitively coloured



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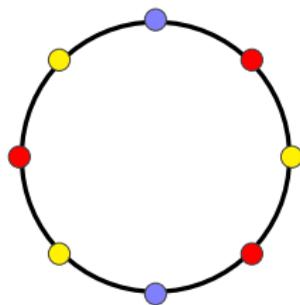


$\pi(G)$  := minimum number of colours in a nonrepetitive colouring of  $G$

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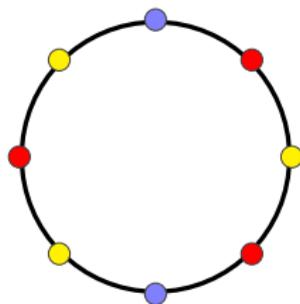
- $\pi(\text{path}) \leq 3$

→ 1906  
[Thue '06]

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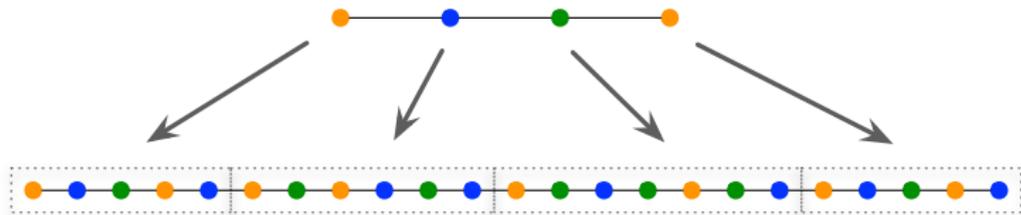
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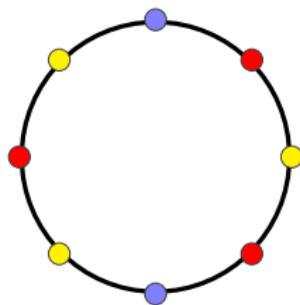
[Thue <sup>1906</sup> '06]



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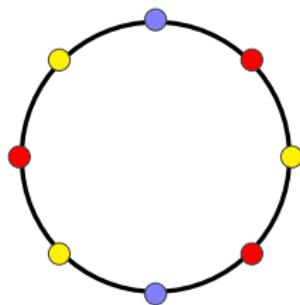
- $\pi(\text{max degree } \Delta \text{ graph}) \leq O(\Delta^2)$

[Alon, Grytczuk, Hałuszczak, Riordan '02]

## nonrepetitive colourings

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[Alon, Grytczuk, Hałuszczak, Riordan '02]

*1906*  
[Thue '06]

open problem [Alon, Grytczuk, Hałuszczak, Riordan '02]

do planar graphs have bounded  $\pi$ ?

Using layered  $H$ -partitions of layered width  $\ell$

- $\pi(H) \leq 4^{\text{tw}(H)}$

[Kündgen & Pelsmayer '08]

## Using layered $H$ -partitions of layered width $\ell$

- $\pi(H) \leq 4^{\text{tw}(H)}$

[Kündgen & Pelsmayer '08]

a colouring is **strongly nonrepetitive** if for every repetitively coloured lazy walk  $v_1, \dots, v_{2t}$ , there exists  $i \in \{1, \dots, t\}$  such that  $v_i = v_{i+t}$

$\pi^*(G) := \min. \#$  colours in a strongly nonrepetitive colouring of  $G$

## Using layered $H$ -partitions of layered width $\ell$

- $\pi(H) \leq 4^{\text{tw}(H)}$

[Kündgen & Pelsmayer '08]

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$\pi^*(G) := \min. \# \text{ colours in a strongly nonrepetitive colouring of } G$

Lemma: [D., Esperet, Joret, Walczak, Wood] 2019

$$\tilde{\pi}^*(G) \leq 4 \cdot \ell \cdot \pi^*(H)$$

Cor:

$$\begin{aligned} \tilde{\pi}(\text{planar } G) &\leq \tilde{\pi}^*(G) \leq 4 \cdot \ell \cdot \pi^*(H) \leq 4 \cdot \ell \cdot 4^{\text{tw}(H)} \\ &\leq 4 \cdot 4^8 = 262,144 \end{aligned}$$

Improved bounds  
and generalizations

## Variant of the main theorem

th: [D., Joret, Micek, Morin, Ueckerdt, Wood] 2019

Every planar graph  $G$  has a layered partition  $\mathcal{P}$  of layered width  $3$  where  $\text{tw}(G/\mathcal{P}) \leq \underline{3}$ .

planar

$H$

best possible

parts not connected

## APPLICATIONS

Coro: [D., Esperet, Joret, Walczak, Wood] 2019

$$\pi(\text{planar } G) \leq 4 \cdot l \cdot 4^{tw(H)} \leq 4 \cdot 3 \cdot 4^3 = 768$$

Coro: [D., Joret, Micoc, Morin, Neckerdt, Wood] 2019

$$\chi_n(\text{planar } G) \leq 3 \cdot l \cdot (2^{tw(H)} - 1) + \frac{3l}{2} \leq 3 \cdot 3 \cdot (2^3 - 1) + 4 \leq 71$$

## APPLICATIONS

Coro: [D., Esperet, Joret, Walczak, Wood] 2019

$$\mathbb{T}(\text{planar } G) \leq 4 \cdot l \cdot 4^{tw(H)} \leq 4 \cdot 3 \cdot 4^3 = 768$$

Coro: [D., Joret, Micoc, Morin, Neckerdt, Wood] 2019

$$\begin{aligned} gn(\text{planar } G) &\leq 3 \cdot l \cdot \left( 2^{tw(H)} - 1 \right) + \lfloor \frac{3l}{2} \rfloor \leq 3 \cdot 3 \cdot (2^3 - 1) + 4 \leq 71 \\ &\leq 3 \cdot l \cdot gn(H) + \lfloor \frac{3}{2} l \rfloor = 3 \cdot 3 \cdot \underline{5} + 4 = \underline{49} \end{aligned}$$

## APPLICATIONS

Cor: [D., Esperet, Joret, Walczak, Wood] 2019  
 $\chi(\text{planar } G) \leq 4 \cdot l \cdot 4^{\text{tw}(H)} \leq 4 \cdot 3 \cdot 4^3 = 768$

Cor: [D., Joret, Micek, Morin, Neckerdt, Wood] 2019  
 $gn(\text{planar } G) \leq 3 \cdot l \cdot (2^{\text{tw}(H)} - 1) + \frac{3l}{2} \leq 3 \cdot 3 \cdot (2^3 - 1) + 4 \leq 71$   
 $\leq 3 \cdot l \cdot gn(H) + \lfloor \frac{3}{2} l \rfloor = 3 \cdot 3 \cdot \underline{5} + 4 = \underline{49}$

Theorem: [Felsner, Micek, Schroeder] 2019+

Planar graphs have  $p$ -centered colourings with  $\Theta(p^3 \lg p)$  col

# Generalization of the main theorem

[D., Joret, Micek, Morin, Ueckerdt, Wood] 2019

Th: Every graph  $G$  of Euler genus  $g$  has a layered partition  $\mathcal{P}$  of layered width  $\max\{2g, 3\}$  where  $\text{tw}(\underbrace{G/\mathcal{P}}_H) \leq 4$ .

→ apex minor free

Th:  $\forall K \exists k, a$  s.t. every  $K$ -minor free graph  $G$  can be obtained by clique-sums of graphs  $G_1, \dots, G_t$  s.t. for  $i \in \{1, \dots, t\}$

$$G_i \subseteq (H_i \boxtimes P_i) + K_a$$

for some  $H_i$  with  $\text{tw}(H_i) \leq k$  and some path  $P_i$

## APPLICATIONS

th: [D., Joret, Micek, Morin, Ueckerdt, Wood] 2019  
graphs excluding a fixed minor have bounded queue #.

th: [D., Esperet, Joret, Walczak, Wood] 2019  
graphs excluding a fixed minor have bounded non- $\chi$

topological

- low tree width colourings

simple proofs  
with better bounds

## OPEN PROBLEMS

- $H$  something other than bounded tree width
- algorithmic and other applications
- proof without using structure theorem
- queue-number, mon-repetitive  $\chi$  of graphs that have structure?  $\leftarrow$  strongly sub-linear separators

the end